

Maximum LPV-SS realization in a static form

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I. FOREWORD

This note is about an augmented state-space realization of LPV input-output forms with static dependency. The importance of this form is due to the fact that it only requires static dependence and it is always realizable. Next step: apply model reduction on it to get to the minimal LPV-SS form with only static dependency.

II. REALIZATION

Consider an LPV-IO representation in the form of

$$y(k) + \sum_{i=1}^{n_a} A_i(p_k)q^{-i}y(k) = \sum_{j=0}^{n_b} B_j(p_k)q^{-j}u(k), \quad (1)$$

where $n_a, n_b \geq 0$, $y : \mathbb{Z} \rightarrow \mathbb{R}^{n_y}$ is the output, $u : \mathbb{Z} \rightarrow \mathbb{R}^{n_u}$ is the input and $p : \mathbb{Z} \rightarrow \mathbb{R}^{n_p}$ is the scheduling signal respectively, while $A_i : \mathbb{R}^{n_p} \rightarrow \mathbb{R}^{n_y \times n_y}$ and $B_i : \mathbb{R}^{n_p} \rightarrow \mathbb{R}^{n_y \times n_u}$ are given matrix functions. The the maximum augmented state-space realization of (1) reads as

$$\left[\begin{array}{c|c} A(p_k) & B(p_k) \\ \hline C(p_k) & D(p_k) \end{array} \right] = \left[\begin{array}{cccccccc|c} -A_1(p_k) & \cdots & -A_{n_a-1}(p_k) & -A_{n_a}(p_k) & B_1(p_k) & \cdots & B_{n_b-1}(p_k) & B_{n_b}(p_k) & B_0(p_k) \\ I & 0 & \cdots & \cdots & \cdots & \cdots & \cdots & 0 & 0 \\ \vdots & \ddots & \ddots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots \\ 0 & \cdots & I & 0 & 0 & 0 & \cdots & 0 & 0 \\ 0 & \cdots & 0 & 0 & 0 & 0 & \cdots & 0 & I \\ 0 & \cdots & 0 & 0 & I & 0 & \cdots & 0 & 0 \\ \vdots & \ddots & \ddots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots \\ 0 & \cdots & \cdots & \cdots & \cdots & 0 & I & 0 & 0 \\ \hline -A_1(p_k) & \cdots & -A_{n_a-1}(p_k) & -A_{n_a}(p_k) & B_1(p_k) & \cdots & B_{n_b-1}(p_k) & B_{n_b}(p_k) & B_0(p_k) \end{array} \right].$$